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Suboptimal Control on Finite Time Interval

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In the paper an algebraic approach to the numerical computation of suboptimal open-loop control on finite time interval is developed. The linear time-invariant continuous system with known transfer functions are considered.

INTRODUCTION

Some optimal control problems can be considered as finite time problems. In this paper only linear time-invariant continuous systems are considered and hence their \mathcal{L} transfer functions are rational.

Our problem can be formulated as that of finding the best approximation g(t) of the given impulse response f(t) in the following sense: the integral

$$\int_0^T (f(t) - g(t))^2 \,\mathrm{d}t$$

is minimized subject to the conditions

 $\begin{aligned} \mathscr{L}(f(t)) &= v(s)/p(s) ,\\ \mathscr{L}(g(t)) &= x(s)/y(s) ,\\ \partial v &< \partial p , \quad \partial x < \partial y , \end{aligned}$

where v, p and x, y are polynomials and ∂v , ∂p and ∂x , ∂y are their respective degrees. For example an approximation of the high order \mathscr{L} transfer function by a low order one may be needed. Assuming the impulse response g(t) partially predetermined by the system with transfer function b/a, $\mathscr{L}(g(t)) = bx/(ay)$, the problem can be described more generally as that of open-loop control. Consider the open-loop control system in Fig. 1, where

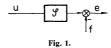
 \mathscr{S} is a realization of the transfer function b(s)/a(s),

u(t) - input signal with \mathscr{L} transform x(s)/y(s),

f(t) - reference signal with \mathscr{L} transform v(s)/p(s),

e(t) – error signal,

then for any given ∂y we can find an input u(t) such that $\int_0^T e^2(t) dt$ is minimized.



Problem formulation:

Let an integer K and polynomials b, a and v, p be given such that $\partial b \leq \partial a$, $\partial v < \partial p$. Find the polynomials x, y such that $\partial y = K$, $\partial x < \partial y$ and the following integral

$$\int_0^T (f(t) - g(t))^2 \,\mathrm{d}t$$

is minimized for

$$f(t) = \mathscr{L}^{-1}\left(\frac{v(s)}{p(s)}\right), \quad g(t) = \mathscr{L}^{-1}\left(\frac{b(s) x(s)}{a(s) y(s)}\right).$$

The solution of this problem is complicated due to finite T.

To begin with we introduce same special notation and operations which are more precisely described in [1] and [2]. The mathematical background is the congruence of analytic functions modulo a polynomial.

Consider a polynomial $m(s) = m_0 + m_1 s + \ldots + m_{\partial m} s^{\partial m}$ with real coefficients and degree $\partial m > 0$. Then the set $\mathcal{M} = \{s : m(s) = 0\}$ is called the spectrum of the polynomial m.

Let functions f, g be analytic on \mathcal{M} . Then f, g are congruent modulo m, written $f = g \mod m$, if there exists a function h analytic on \mathcal{M} such that f = g + hm. For any function f analytic on \mathcal{M} only one polynomial r exists such that

(1)
$$f = r \mod m, \quad \partial r < \partial m.$$

The polynomial m is called the modulus. The operation which yields such a polynomial r is called the reduction of f modulo m and it is denoted as

$$[f]_m = r \,.$$

Denote \mathscr{F}_m the set of all functions analytic or having at worst removable singularities on \mathscr{M} .

PROPERTIES OF REDUCTION MODULO m

Let a modulus m, $\partial m > 0$, and function $f, g \in \mathscr{F}_m$ be given. Then for $[f]_m = a$, $[g]_m = b$ and any complex number λ the next equations hold:

(3)
$$[f+g]_m = [f]_m + [g]_m = a + b$$
,

(4)
$$[\lambda f]_m = \lambda [f]_m = \lambda a ,$$

(5)
$$[fg]_m = [[f]_m [g]_m]_m = [ab]_m,$$

(6)) if
$$f/g \in \mathscr{F}_m$$
 then

$$\left[\frac{f}{g}\right]_m = \left[\frac{[f]_m}{[g]_m}\right]_m = \left[\frac{a}{b}\right]_m.$$

If the function f is a polynomial then the reduction of f modulo m produces the remainder after dividing f by m. Procedures for the computation of $[f]_m$ for the functions $\ln(s)$, e^{ks} , \sqrt{s} , s^k with k real, and for b/a, $b \cdot a$ with polynomials a, b are described in [1].

Now we introduce a new operation

(7)
$$\left\langle \frac{s[f]_a}{a} \right\rangle_s = \lim_{s \to \infty} \frac{s[f]_a}{a} = \frac{c_{k-1}}{a_k},$$

where $k = \partial a$

$$[f]_a = c_0 + c_1 s + \ldots + c_{k-1} s^{k-1}$$

and the subscript s denotes the variable with respect to which the operation $\langle . \rangle_s$ is performed.

The next Theorem is proved in [2].

Theorem 1. Let polynomials b, a, $\partial b < \partial a$, with real coefficients be given. Then the inverse Laplace transform of b/a denoted as $\mathscr{L}^{-1}(b/a)$, is the real-valued function f(t) given by

(8)
$$f(t) = \left\langle \frac{s}{a} \left[b \ e^{st} \right]_a \right\rangle_a$$

and furthermore

$$\frac{\mathrm{d}^{i}}{\mathrm{d}t^{i}}f(t) = \left\langle \frac{s}{a} \left[s^{i}b \ \mathrm{e}^{st} \right]_{a} \right\rangle_{s}, \quad i = 1, 2, \ldots$$

EVALUATION $\int_0^T f(t) g(t) dt$

The basic formula of this paper is in the following Theorem.

Theorem 2. Let polynomials b, a, v, p, $\partial b < \partial a$, $\partial v < \partial p$ be given. Denote $f(t) = \mathscr{L}^{-1}(b/a)$, $g(t) = \mathscr{L}^{-1}(v/p)$ then

(9)
$$I = \int_{0}^{T} f(t) g(t) dt = \left\langle \frac{s}{a} \left[b \frac{\overline{v} - e^{sT} [\overline{v} e^{-sT}] \overline{p}}{\overline{p}} \right]_{a} \right\rangle_{s} = \left\langle \frac{s}{p} \left[v \frac{\overline{b} - e^{-sT} [\overline{b} e^{-sT}] \overline{a}}{\overline{a}} \right]_{p} \right\rangle_{s},$$

where $\bar{v}(s) = v(-s)$, $\bar{b}(s) = b(-s)$.

Proof. Using Theorem 1 and the linearity of the operation $\langle . \rangle_s$ we can write

$$I = \left\langle \frac{s}{a} \left[b \int_0^T e^{st} g(t) dt \right]_a \right\rangle_s.$$

It is evident that

$$\int_0^T e^{st} g(t) dt = \int_0^\infty e^{st} g(t) dt - \int_T^\infty e^{st} g(t) dt$$

and $\int_0^\infty e^{st} g(t) dt = \bar{v}/\bar{p}$ by definition of \mathscr{L} transform. By Theorem 1

$$g(t) = \left\langle \frac{s}{p} \left[v \, \mathrm{e}^{st} \right]_p \right\rangle_s$$

and

$$g(T + \tau) = \left\langle \frac{s}{p} \left[v \, \mathrm{e}^{sT} \, \mathrm{e}^{s\tau} \right]_{p} \right\rangle_{s} = \left\langle \frac{s}{p} \left[\left[v \, \mathrm{e}^{sT} \right]_{p} \, \mathrm{e}^{s\tau} \right]_{p} \right\rangle_{s}.$$

Hence

$$\int_{T}^{\infty} e^{st} g(t) dt = e^{sT} \int_{0}^{\infty} e^{s\tau} g(T+\tau) d\tau = e^{sT} \frac{\left[\overline{v} e^{-sT}\right]_{\overline{p}}}{\overline{p}}$$

and the proof is complete. Considering $f(t) = \mathscr{L}^{-1}(v|p), g(t) = \mathscr{L}^{-1}(b|a)$ the second formula follows.

Remark. The computation of the formula (9) can be rearranged in the following way. Denote

(10)
$$\left[\bar{v} e^{-sT}\right]_{\bar{p}a} = x ,$$

then $\partial x < \partial p + \partial a$. Use division algorithm for x/\bar{p} , then

(11)
$$x = \bar{p}q + r, \quad \partial r < \partial \bar{p}$$

and

$$I = \int_0^T f(t) g(t) dt = \left\langle \frac{s}{a} \left[b e^{sT} \frac{x - [x]_{\bar{p}}}{\bar{p}} \right]_a \right\rangle_s = \left\langle \frac{s}{a} \left[b e^{sT} q \right]_a \right\rangle_s.$$

MINIMIZATION PROCEDURE

Denote

$$g(t) = \mathscr{L}^{-1}\left(\frac{bx}{ay}\right), \quad f(t) = \mathscr{L}^{-1}\left(\frac{v}{p}\right),$$

and

$$I = \int_0^T (f(t) - g(t))^2 \, \mathrm{d}t \, .$$

Assuming the polynomials a, y, p in normalized form, $a_{\hat{c}a} = 1$, $y_{\hat{c}y} = 1$, $p_{\hat{c}p} = 1$, we can write

$$f(t) = \left\langle \frac{s}{s^{e_p}} \left[e^{st} v \right]_p \right\rangle_s,$$
$$g(t) = \left\langle \frac{s}{s^{e_a + e_y}} \left[e^{st} b x \right]_{ay} \right\rangle_s.$$

From the problem formulation the next formula follows.

(12)
$$I(x, y) = \int_0^T (f(t) - g(t))^2 dt.$$

Using the variation approach the condition for minimizing (12) is given by

(13)
$$\delta I = 2 \int_0^T (f(t) - g(t)) \, \delta g(t) \, \mathrm{d}t = 0$$

for any variation of g(t) which is given as

$$\delta g(t) = \mathscr{L}^{-1} \left(\frac{b \delta x}{a y} - \frac{b x}{a y^2} \delta y \right).$$

Hence using Theorem 2

(14)
$$\delta I = \left\langle \frac{1}{s^{\hat{c}a+\hat{c}y}} \left[b\delta x(\bar{F}-\bar{G}) \right]_{ny} \right\rangle_{s} - \left\langle \frac{1}{s^{\hat{c}a+\hat{c}y}} \left[bx\delta y(\bar{F}-\bar{G}) \right]_{ny^{2}} \right\rangle_{s} = 0,$$

for all δx , δy , where

$$\overline{F} = \frac{\overline{v} - e^{+sT} [e^{-sT} \overline{v}]_{\overline{p}}}{\overline{p}},$$
$$\overline{G} = \frac{\overline{b}\overline{x} - e^{sT} [e^{-sT} \overline{b}\overline{x}]_{\overline{a}\overline{y}}}{\overline{a}\overline{y}},$$
$$\overline{a}(s) = a(-s).$$

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$$\begin{split} \delta x &= \delta x_0 + s \delta x_1 + \ldots + s^{\delta y - 1} \delta x_{\delta y - 1} ,\\ \delta y &= \delta y_0 + s \delta y_1 + \ldots + s^{\delta y - 1} \delta y_{\delta y - 1} ,\\ y_{\delta y} &= 1 \end{split}$$

the equation (14) may be written as a set of conditions

(15)
$$\frac{\partial I}{\partial x_i} = I_1^{(i)} = \left\langle \frac{s}{s^{\delta a + \delta y}} \left[s^i b(\bar{F} - \bar{G}) \right]_{ay} \right\rangle_{s} = 0,$$
$$\frac{\partial I}{\partial y_i} = I_2^{(i)} = \left\langle \frac{s}{s^{\delta a + 2\delta y}} \left[s^i b x(\bar{F} - \bar{G}) \right]_{ay2} \right\rangle_{s} = 0,$$
$$i = 0, 1, \dots, \partial y - 1.$$

To solve these nonlinear equations we use Newton's iteration method in the form

(16)
$$I_1^{(i)} + \delta I_1^{(i)} = 0$$
, for $i = 0, 1, ..., \partial y - 1$,
 $I_2^{(i)} + \delta I_2^{(i)} = 0$,

where $\delta I_1^{(i)}$, $\delta I_2^{(i)}$ denote variations of $I_1^{(i)}$, $I_2^{(i)}$ with respect to δx , δy .

Note that the application of Theorem 2 to equations (15) gives

$$I_1^{(i)} = \int_0^T (f(t) - g(t)) h^{(i)}(t) dt = 0$$

where $h^{(i)} = \mathcal{L}^{-1}(s^{i}H)$, H = b/ayand similarly

$$I_2^{(i)} = \int_0^T (f(t) - g(t)) c^{(i)}(t) dt = 0$$

where

$$c^{(i)}(t) = \mathscr{L}^{-1}(s^{i}C), \quad C = \frac{bx}{ay^{2}}.$$

Hence

$$\delta I_1^{(i)} = \int_0^T ((f(t) - g(t)) \,\delta \,h^{(i)}(t) - h^{(i)}(t) \,\delta \,g(t)) \,\mathrm{d}t \,,$$

$$\delta I_2^{(i)} = \int_0^T ((f(t) - g(t)) \,\delta c^{(i)}(t) - c^{(i)}(t) \,\delta \,g(t)) \,\mathrm{d}t \,,$$

where

$$\delta h^{(i)}(t) = \mathscr{L}^{-1}\left(-\frac{s^i b}{a y^2} \delta y\right),$$

$$\delta c^{(i)}(t) = \mathscr{L}^{-1}\left(\frac{s^i b \delta x}{a y^2} - 2\frac{s^i b x \delta y}{a y^3}\right),$$

$$\delta g(t) = \mathscr{L}^{-1}\left(\frac{b \delta x}{a y} - \frac{b x \delta y}{a y^2}\right).$$

Define

$$\begin{aligned} \overline{H}^{(i)} &= \frac{\overline{b}(-s)^i - \mathrm{e}^{sT}[\mathrm{e}^{-sT}(-s)^i \,\overline{b}]_{\bar{a}\bar{y}}}{\bar{a}\bar{y}}, \\ \overline{C}^{(i)} &= \frac{\overline{b}\overline{x}(-s)^i - \mathrm{e}^{sT}[\mathrm{e}^{-sT}b\overline{x}(-s)^i]_{\bar{a}\bar{y}\bar{p}^2}}{\bar{a}\bar{y}^2} \end{aligned}$$

and the elements of the matrices A, B, C, D of the dimension $\partial y \times \partial y$ in the following way.

(17)
$$A_{i+1,j+1} = \left\langle \frac{s}{s^{\partial a+\partial y}} \left[b s^{j} \overline{H}^{(i)} \right]_{ay} \right\rangle_{s},$$

$$B_{i+1,j+1} = \left\langle \frac{s}{s^{\partial a+2\partial y}} \left[bxs^{j}\overline{H}^{(i)} \right]_{ay^{2}} \right\rangle_{s} - \left\langle \frac{s}{s^{\partial a+2\partial y}} \left[bs^{i+j}(\overline{F}-\overline{G}) \right]_{ay^{2}} \right\rangle_{s},$$

$$C_{i+1,j+1} = B_{j+1,i+1},$$

$$D_{i+1,j+1} = \left\langle \frac{s}{s^{\partial a+3\partial y}} \left[2bs^{i+j}(\overline{F}-\overline{G}) \right]_{ay^{3}} \right\rangle_{s} - \left\langle \frac{s}{s^{\partial a+2\partial y}} \left[bxs^{j}\overline{C}^{(i)} \right]_{ay^{2}} \right\rangle_{s},$$

$$i = 0, 1, \dots, \partial y - 1, \quad j = 0, 1, \dots, \partial y - 1.$$

Using Theorem 2 it is simple to prove that A, D are symmetric matrices. Using the matrix notation

$$J_{1}(X, Y) = \begin{bmatrix} I_{1}^{(0)} \\ I_{1}^{(1)} \\ \vdots \\ I_{1}^{(2y-1)} \end{bmatrix}, \quad J_{2}(X, Y) = \begin{bmatrix} I_{2}^{(0)} \\ I_{2}^{(1)} \\ \vdots \\ I_{2}^{(2y-1)} \end{bmatrix},$$
$$X = \begin{bmatrix} x_{0} \\ x_{1} \\ \vdots \\ x_{\partial y-1} \end{bmatrix}, \quad Y = \begin{bmatrix} y_{0} \\ y_{1} \\ \vdots \\ y_{\partial y-1} \end{bmatrix}, \quad y_{\partial y} = 1,$$

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(18)
$$\delta J_1 = -A\delta X + B\delta Y,$$
$$\delta J_2 = B'\delta X + D\delta Y.$$

Let us point out that $I_1^{(i)}$ can be computed as

$$I_{1}^{(i)} = \left\langle \frac{s}{s^{\delta a + \delta y}} \left[s^{i} b \overline{F} \right]_{ay} \right\rangle_{s} - \left\langle \frac{s}{s^{\delta a + \delta y}} \left[b x \overline{H}^{(i)} \right]_{ay} \right\rangle_{s}$$
$$I_{1}^{(i)} = \left\langle \frac{s}{s^{\delta a + \delta y}} \left[s^{i} b \overline{F} \right]_{ay} \right\rangle_{s} - \sum_{j=0}^{\delta y-1} A_{i+1,j+1} x_{j}$$

and $I_2^{(1)}$ corresponds to the second term in the formula for $B_{i+1,1}$. Hence in our iteration process, we can define for given Y the optimal X by

(19)
$$X = A^{-1} J_1(0, Y).$$

Substituting (19) into (18) and using Newton's formula

$$-A\delta X \times B\delta Y = 0,$$

$$J_2 + B'\delta X + D\delta Y = 0$$

we obtain

$$\delta Y = -(D + B'A^{-1}B)^{-1} J_2 \, .$$

Now we summarize the iteration algorithm.

- (1) Start from initial condition $y = s^{\partial y}$.
- (2) Compute optimal X given by $X = A^{-1}J_1$.
- (3) Compute the value $Y = Y + \delta Y$, where

$$\delta Y = -(D + B'A^{-1}B)^{-1} J_2.$$

- (4) If the *n*-th iteration and the (n-1)-th iteration are not sufficiently near go to (2).
- (5) Print results.

This algorithm was programed and tested on IBM 370/135 computer using PL/I language.

Example 1. Consider the reference signal $f(t) = \mathscr{L}^{-1}(v/p)$, where

$$v = 0.03387 - 0.06093s - 1.1634s^2 - 0.43164s^3 + 8.6775s^4 - - 9.525s^5 + 0.9898s^6 + 1.5379s^7,$$

$$p = 0.01129 - 0.03329s - 1.1056s^2 + 5.82213s^3 - 1.77933s^4 - - 6.36952s^5 + 1.97685s^6 + s^7 + 0.5s^8,$$

and the system transfer function b/a, where

$$b = 1$$
,
 $a = 1 + 0.2s + 0.04s^2$.

Find polynomials x, y such that $\partial y = 3$ and the integral

$$\int_0^5 (f(t) - g(t))^2 \, \mathrm{d}t \,, \quad \text{where} \quad g(t) = \mathscr{L}^{-1}\left(\frac{bx}{ay}\right),$$

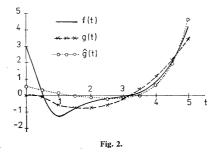
is minimized.

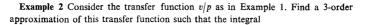
The result is obtained after five iterations in the form

$$x = 0.8023 - 1.225s + 0.1849s^{2},$$

$$y = -0.01124 - 0.007296s - 0.000608s^{2} + s^{3}.$$

The functions f(t), g(t) are plotted in Fig. 2.





$$\int_{0}^{5} (f(t) - \hat{g}(t))^{2} dt, \text{ where } \hat{g}(t) = \mathscr{L}^{-1}(x/y)$$

is minimized. The solution is similar as in the previous Example, but b = 1, a = 1.

After five iterations we obtain the result

$$x = 0.69452 - 1.2376s + 0.5877s^2,$$

$$y = 0.6626 - 0.6487s - 1.011s^2 + s^3.$$

The function $\hat{g}(t)$ is plotted in Fig. 2.

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