

# Convergence Properties of Adaptive Threshold Elements in Respect to Application and Implementation

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Referring to the convergence proof of the error correction algorithm estimations of the upper bound  $t_g$  of correction steps and of the sufficient weight interval  $(-w_g, +w_g)$  in digital implementation are derived and illustrated. Both  $t_g$  and  $w_g$  depend on the gap  $2\hat{d}$  between the convex hulls of classes in pattern space. The estimations give  $t_g \sim |x|_{\max}^2/\hat{d}^2$  and  $w_g \sim |x|_{\max}^2/\hat{d}$ , where  $|x|_{\max}$  is the largest pattern vector of both classes.

## 1. INTRODUCTION

Adaptive threshold elements (ATE) are used to realize or to modify linear discriminant functions in a training (learning) process. ATE are applied as basic building blocks for learning systems (hardware) and for iterative determination of discriminant functions with aid of computers (software). They are simple realizable. There exist numerous detailed descriptions of the convergence properties of training algorithms and of the classification properties, e.g. [1, 2, 5, 6 and 9].

The convergence proofs show that for the case of linear separability a solution is found after a finite number of correction steps. For practical purposes it is unsatisfactory proving only the finiteness of the number of correction steps without giving a bound. In the third part of this paper an estimation of the upper bound in relation to the gap between the convex hulls of pattern classes is derived.

The fast progress of digital circuits has an important influence on the recent implementations of ATE, e.g. [3, 13 and 11]. These constructions are characterized by a limited number of bits for the parameter values of the discriminant function. Therefore the possibilities for the implementation are restricted. In the fourth part a relation is derived between the gap width and the weight interval ensuring convergence. The results of the 3rd and 4th part are connected in the 5th part.

For illustration the estimations of  $t_g$  and  $w_g$  are applied on the special case of threshold functions in BOOLEAN space.

In the following part the properties of ATE are illustrated and basic definitions are introduced.

## 2. PATTERN SPACE AND $\delta$ -SEPARABILITY

Patterns are characterized by  $n$ -component vectors  $x$  in  $\mathcal{X}$ -space. Let  $\mathcal{K}1$  and  $\mathcal{K}2$  be two pattern classes. In the linearly-separable case a pair of hyperplanes exists separating these classes. The hyperplanes are determined by the normal vector  $w^+ / |w^+|$ , the signed distance  $a$  from origin and the distance  $2\Delta > 0$  between the hyperplanes (Fig. 1).\*)  $2\hat{\Delta}$  denotes the distance between the convex hulls of  $\mathcal{K}1$  and  $\mathcal{K}2$ .

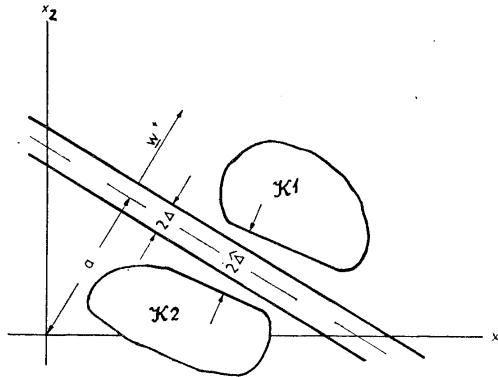


Fig. 1. Two-class problem with hyperplane distance  $2\Delta$  and gap width  $2\hat{\Delta}$  in  $\mathcal{X}$ -space.

The decision rule is

$$(1) \quad \frac{w^+}{|w^+|} x - a \begin{cases} \geq \Delta & \Rightarrow x \in \mathcal{K}1 \\ \leq -\Delta & \Rightarrow x \in \mathcal{K}2 \\ \text{otherwise} & \Rightarrow \text{rejection.} \end{cases}$$

The inner product of vectors is not especially indicated. For the following considerations it is useful to introduce  $\mathcal{Y}$ -space (defined by the augmented pattern vector  $y$ ), the weight vector  $w$  and a parameter  $\delta$ :

\*) Notice that in all figures underlined letters denote vectors (printed by bold letters in the text).

$$(2) \quad \begin{aligned} w &= (w^+, w_{n+1}) \\ y &= (x, x_{n+1}) \quad \text{for } x \in \mathcal{K}1 \\ y &= (-x, -x_{n+1}) \quad \text{for } x \in \mathcal{K}2 \\ \delta &> 0. \end{aligned}$$

Using Eq. (2) we can write for Eq. (1)

$$(3) \quad wy \geq \delta,$$

$$(4) \quad \text{with } a = -x_{n+1}w_{n+1}/|w^+| \text{ and } A = \delta/|w^+|.$$

To separate  $\mathcal{K}1$  and  $\mathcal{K}2$  we have to find a vector  $w$  satisfying  $wy > 0$  for all  $y$ . One of the algorithms doing this is the following well known error correction algorithm

$$(5) \quad \begin{aligned} w_1 &\text{ arbitrary} \\ w_{t+1} &= w_t + \gamma y_t \quad \text{if } w_t y_t \leq 0 \\ &\text{with } 0 < \varepsilon \leq \gamma \leq c < \infty, \end{aligned}$$

where  $\gamma$  is the correction factor and  $t$  denotes the correction step. Only patterns  $y$  misclassified during the adaption process cause corrections and are accounted to the  $y_t$  sequence.

Bounds for the number of correction steps and the sufficient weight interval can be obtained only when some assumptions are made concerning the type of separability. Starting point is a definition of separability with gap. A definition given in [10] will be expanded.

**Definition 1.** Two classes  $\mathcal{K}1$  and  $\mathcal{K}2$  are said to be  $\delta_1$ -separable if there exists a weight vector  $w$  with

$$wy \geq \delta_1 > 0 \quad \text{and} \quad |w_i| \leq 1, \quad i = 1, 2, \dots, n+1.$$

**Definition 2.** Two classes  $\mathcal{K}1$  and  $\mathcal{K}2$  are said to be  $\max\text{-}\delta_1$ -separable if they are  $\delta_1$ -separable but not  $(\delta + \varepsilon)_1$ -separable with  $\varepsilon > 0$ . Figure 2 shows examples for the Definitions 1 and 2.

### 3. UPPER BOUND OF CORRECTION STEPS

Linear separability is assumed. Then the proof shows that algorithm (5) yields a solution vector after a limited number of correction steps, and the correction bound is especial a function of the gap  $2\hat{A}$ .

Let  $w_\delta$  be a solution vector of the region  $\mathcal{W}_\delta$  containing all  $w$  with  $wy \geq \delta$  (Fig. 3). Then each  $\alpha w_\delta$  for  $\alpha > 0$  is a solution vector of the solution region  $\mathcal{W}_0$  containing all  $w$  with  $wy > 0$  (Fig. 3). We only consider such  $y$  of the (cyclical) pattern sequence

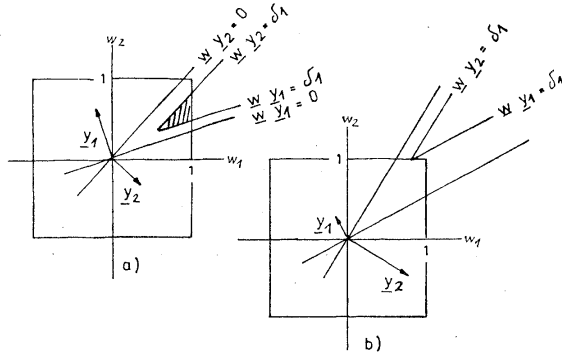


Fig. 2. Illustration of the Definitions 1 and 2 in the weight space. Each class contains only one pattern. The hatched region for  $\delta_1$ -separability in a) degenerates to one point for  $\max\text{-}\delta_1$ -separability in b).

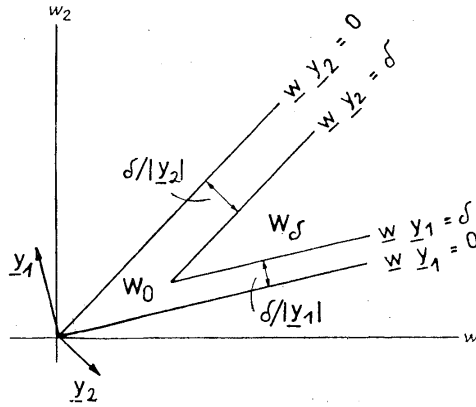


Fig. 3. Solution region  $W_0$  for  $w y > 0$  and  $W_\delta$  for  $w y \geq \delta$ .

which lead to a correction. Then the following relations are valid (compare also [2]):

$$\begin{aligned}
 w_{t+1} &= w_t + \gamma y_t, \\
 w_{t+1} - \alpha w_\delta &= w_t - \alpha w_\delta + \gamma y_t, \\
 |w_{t+1} - \alpha w_\delta|^2 &= |w_t - \alpha w_\delta|^2 + 2\gamma w_t y_t - 2\gamma \alpha w_\delta y_t + \gamma^2 |y_t|^2,
 \end{aligned}
 \tag{6}$$

with  $w_t y_t \leq 0$ ,  $w_\delta y_t \geq \delta$  and  $|y_t| \leq |y|_{\max}^2$  Eq. (6) yields

$$(7) \quad |w_{t+1} - \alpha w_\delta|^2 \leq |w_t - \alpha w_\delta|^2 - 2\gamma\alpha\delta + \gamma^2 |y|_{\max}^2.$$

Starting with the weight vector  $w_1$  we obtain after  $t$  corrections

$$(8) \quad |w_{t+1} - \alpha w_\delta|^2 \leq |w_1 - \alpha w_\delta|^2 - t(2\gamma\alpha\delta - \gamma^2 |y|_{\max}^2).$$

If  $w_1 = 0$  is chosen and  $0 \leq |w_{t+1} - \alpha w_\delta|$  is taken into consideration, then from Eq. (8) follows

$$0 \leq \alpha^2 |w_\delta|^2 - t(2\gamma\alpha\delta - \gamma^2 |y|_{\max}^2).$$

After rearrangement we obtain

$$(9) \quad t \leq \frac{\alpha^2 |w_\delta|^2}{2\gamma\alpha\delta - \gamma^2 |y|_{\max}^2}.$$

Since  $t$  must be positive the same follows for the denominator of Eq. (9) and finally

$$(10) \quad \alpha > \frac{\gamma |y|_{\max}^2}{2\delta}.$$

Relation (10) ensures the step by step approach of  $w_t$  to  $\alpha w_\delta$ . With  $t_g$  as upper bound for the number of correction steps in Eq. (9) the minimization in respect to  $\alpha$  yields for

$$(11) \quad \alpha = \frac{\gamma |y|_{\max}^2}{\delta}$$

the correction bound

$$(12) \quad t_g = \frac{|y|_{\max}^2 |w_\delta|^2}{\delta^2}$$

in  $\mathcal{Y}$ -space. The estimation of correction bound  $t_g$  in Eq. (12) gives a minimum if, for a given  $\delta$ , the shortest solution vector is taken as  $w_\delta$  or, for  $\max$ - $\delta_1$ -separability,  $\delta_1$  and  $\hat{w}(\delta_1)$  are put in.

However, a relation between the correction bound  $t_g$  and the gap  $2\Delta$  in  $\mathcal{X}$ -space would be quite clearer. Eq. (12) can be written as

$$(13) \quad t_g = (|x|_{\max}^2 + x_{n+1}^2) \frac{|w^+|^2 + w_{n+1}^2}{\delta^2}.$$

With  $\Delta = \delta/|w^+|$  and  $a = -w_{n+1}x_{n+1}/|w^+|$  follows

$$t_g = (|x|_{\max}^2 + x_{n+1}^2) \frac{x_{n+1}^2 + a^2}{x_{n+1}^2 \Delta^2}.$$

164 We only know  $|a| < |x|_{\max}$  about the value of  $a$  and with  $\hat{\Delta}$  for  $\Delta$  we obtain

$$(14) \quad t_g = \frac{(|x|_{\max}^2 + x_{n+1}^2)^2}{x_{n+1}^2 \hat{\Delta}^2}.$$

Figure 4 illustrates Eq. (14).

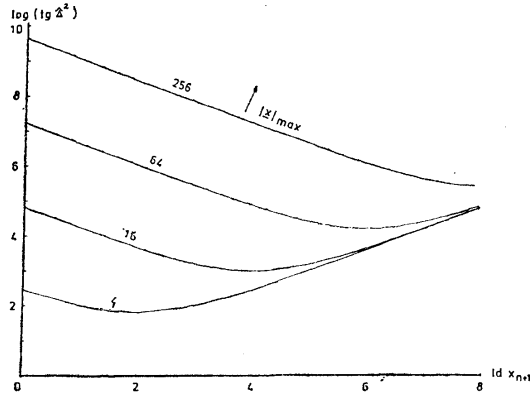


Fig. 4.  $\log(t_g \hat{\Delta}^2)$  as function of  $\text{ld}(x_{n+1})$  with  $|x|_{\max}$  as parameter.

Figure 4 shows that an inadequate choice of the  $(n+1)$ th component can increase  $t_g$  at some orders. In the literature  $x_{n+1}$  is in general chosen independently on  $|x|_{\max}$ . The minimum of  $t_g$  with respect to  $x_{n+1}$  is

$$(15) \quad t_g = 4 \frac{|x|_{\max}^2}{\hat{\Delta}^2} \quad \text{for } x_{n+1} = |x|_{\max}.$$

Eq. (15) gives a direct dependence of the correction bound  $t_g$  on the gap  $2\hat{\Delta}$  and the largest pattern vector  $|x|_{\max}$  in  $\mathcal{X}$ -space.

The rearrangement of Eq. (15) results in

$$(16) \quad \hat{\Delta} \leq 2|x|_{\max}/\sqrt{t}.$$

After  $t$  corrections the following statement can be made, the gap is less than or equal to the right side of Eq. (16).

Theorem 1 gives the basis for the convergence of algorithm (5) in case of a limited number of bits for each weight.

**Theorem 1.** If  $\mathcal{H}1$  and  $\mathcal{H}2$  are  $\delta_1$ -separable and if  $\gamma$  is chosen

$$(17) \quad 0 < \gamma < \frac{2\delta_1}{|y|_{\max}^2}$$

then algorithm (5) with the additional condition  $|w_i| \leq 1$ ,  $i = 1, 2, \dots, n+1$ , leads to a solution after a finite number of corrections.

For proving this we go back to part 3. From Eqs. (7) resp. (10) follows that for

$$0 < \gamma < \frac{2\delta}{|y|_{\max}^2}$$

$w$  approaches any  $w \in \tilde{\mathcal{W}}_\delta$  at each correction. For proving this put  $\alpha = 1$  and determine  $\gamma$  for

$$\gamma \text{ to } |w_{t+1} - w_\delta|^2 - |w_t - w_\delta|^2 < 0.$$

How is the influence of weight limitation? Let  $w_a$  be a vector outside the bounded hypercube and  $w_b$  the nearest vector to  $w_a$  inside the bounded hypercube, then we can show for all vectors  $w$  inside

$$|w_b - w| < |w_a - w|,$$

particularly for  $w \in \tilde{\mathcal{W}}_\delta$ , too, where  $\tilde{\mathcal{W}}_\delta$  is the bounded region of  $\mathcal{W}_\delta$  (hatched in Fig. 2). The vector  $w$  enters the region  $\mathcal{W}_0$  during the step by step approach of  $w$  to  $\tilde{\mathcal{W}}_\delta$ . This completes the proof.

For

$$\gamma > \frac{2\delta_1}{|y|_{\max}^2 (1 - \delta_1^2)}$$

examples can be given, where  $w$  oscillates during training (Figure 5). Therefore estimation (17) can not be essentially improved.

For digital implementations (pattern components and weighting coefficients are digital) algorithm (5) has to be chosen with integer  $\gamma$ . An appropriate choice is  $\gamma = \tilde{\gamma} = 1$ . From Theorem 1 follows:

**Corollary 1.** Supposing  $\delta_1$ -separability and putting  $\gamma = \tilde{\gamma} = 1$  in algorithm (5) the following interval  $(-w_g, +w_g)$  of weights is sufficient for convergence:

$$(18) \quad w_g > \frac{|y|_{\max}^2}{2\delta_1}.$$





resp.

$$\Theta > \frac{x_{n+1} \hat{d}}{|y|_{\max}}.$$

In case of  $\max\text{-}\delta_1\text{-separability}$  there is at least one  $|w_i| = 1$  and hence  $|w| \geq 1$  as well as

$$\delta_1 \geq \Theta > x_{n+1} \hat{d} / |y|_{\max}.$$

Inserted in Eq. (18) and with  $\hat{d}$  for  $d$  we obtain

$$w_\theta = \frac{|y|_{\max}^3}{2x_{n+1} \hat{d}}$$

resp.

$$(21) \quad w_\theta = \frac{(x_{\max}^2 + x_{n+1}^2)^{3/2}}{2x_{n+1} \hat{d}}.$$

Figure 6 illustrates the influence of  $x_{n+1}$  on the sufficient interval of weights. The

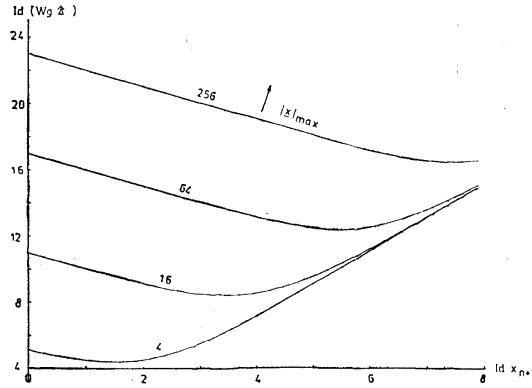


Fig. 6.  $\text{ld}(w_\theta \hat{d})$  as function of  $\text{ld}(x_{n+1})$  with  $|x|_{\max}$  as parameter.

minimum of Eq. (21) with respect to  $x_{n+1}$  is

$$(22) \quad w_\theta = \frac{\sqrt{27} |x|_{\max}^2}{4 \hat{d}} \quad \text{for} \quad x_{n+1} = |x|_{\max} / \sqrt{2}.$$

168 If  $w_g$  is fixed by the implementation then the convergence of algorithm (5) is ensured if

$$\hat{\Delta} \geq \frac{\sqrt{27}}{4} \frac{|x|_{\max}^2}{w_g} \quad \text{and} \quad x_{n+1} = |x|_{\max}/\sqrt{2}.$$

Demanding the separation of classes with a zone width  $2\tilde{\Delta}$  near  $2\hat{\Delta}$  Eq. (19) yields for the sufficient weight interval

$$(23) \quad w_g = \frac{(x_{\max}^2 + x_{n+1}^2)^{3/2}}{2x_{n+1}(\hat{\Delta} - \tilde{\Delta})} \quad \text{for} \quad 0 \leq \tilde{\Delta} < \hat{\Delta}$$

resp.

$$(24) \quad w_g = \frac{\sqrt{27}}{4} \frac{|x|_{\max}^2}{\hat{\Delta} - \tilde{\Delta}} \quad \text{for} \quad 0 \leq \tilde{\Delta} < \hat{\Delta} \quad \text{and} \quad x_{n+1} = |x|_{\max}/\sqrt{2}.$$

Figure 7 shows the hyperbolic growth of  $w_g$  for  $\tilde{\Delta} \rightarrow \hat{\Delta}$ .

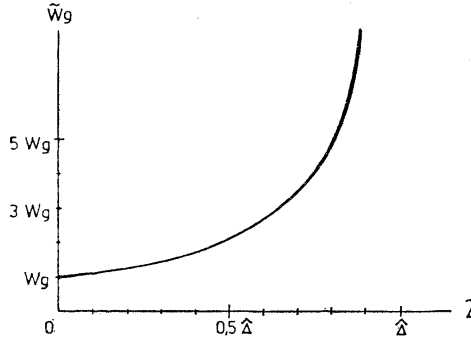


Fig. 7. Sufficient interval of weights  $(-w_g, +w_g)$  depending on  $\hat{\Delta}$  and  $\tilde{\Delta}$ .

## 5. INTERVAL OF WEIGHTS AND CORRECTION BOUND

In what respect does the interval of weights influence the statements for the correction bound? The basis for estimating the correction bound  $t_g$  is the step by step approach of  $w$  to a vector  $\alpha w_g$  with  $w_g \in \mathcal{W}_g$ . Is  $w_g$  specified by  $\hat{w}$ , where  $\hat{w}$  is the solution vector for *max- $\delta_1$ -separability*, Eq. (11) yields for  $\gamma = 1$  an  $\alpha = |y|_{\max}^2/\delta_1$ . As in Figure 8 is shown,  $\alpha\hat{w}$  lies in a hypercube with the half edge

$$(25) \quad w_\alpha = \frac{|y|_{\max}^2}{\delta_1}.$$

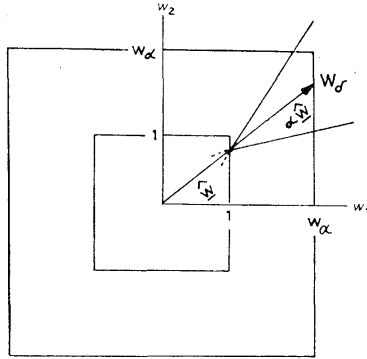


Fig. 8. Hypercube with  $\hat{w}$  for  $\max\text{-}\delta_1$ -separability and  $\alpha\hat{w}$  of the convergence proof to algorithm (5).

By a limitation of weight to  $|w_i| \leq |y|_{\max}^2 / \delta_1$  the correction bound  $t_g$  according to Eq. (12) is not increased since the impact of limitation at a correction step only leads to a stronger approach of  $w$  to  $\alpha\hat{w}$  in this step (compare the proof of Theorem 1).

The comparison of  $w_\alpha$  from Eq. (25) and  $w_g$  from Eq. (18) leads to

$$(26) \quad w_\alpha = 2w_g.$$

The statements for the correction bound in part 4 allow a limited interval for the weights of  $\pm w_\alpha = \pm 2w_g$ .

## 6. EXAMPLE: BOOLEAN THRESHOLD FUNCTIONS

The separation problem in **BOOLEAN** space is a well known special case. For lower dimensionality the interval  $(-w_{\max}, +w_{\max})$  of weights is known being sufficient for the realization of all threshold functions according to

$$wy \geq 1 \quad \text{with} \quad y_i = \pm 1$$

under the condition

$$\sum_{i=1}^{n+1} |w_i| = \text{minimum}$$

[4, 7 and 12]. Therefore we have all information to determine  $t_g$  and  $w_g$ .

The number of all existing threshold functions and the maximum weight  $w_{\max}$  for all

170 dimensions  $\leq 8$  are given in the 2nd and 3rd row of Table 1. For the dimension  $n$  we obtain  $|y|_{\max}^2 = n + 1$ . With  $\delta_1 = 1/|w|_{\max}$  Eqs. (12) and (18) lead to

$$(27) \quad t_g = (n + 1)^2 w_{\max}^2$$

and

$$(28) \quad w_g = \frac{n + 1}{2} w_{\max}.$$

Table 1

dimension $n$	1	2	3	4	5	6	7	8
threshold function	4	14	104	1882	94572	15028134	$8,38 \cdot 10^9$	$1,76 \cdot 10^{13}$
maximum weight $w_{\max}$	1	1	2	3	5	9	18	42
correction bound $t_g$	4	9	32	225	900	3969	20736	142884
sufficient weight $w_g$	1	1,5	4	7,5	15	31,5	72	189
bits for $(-w_g, +w_g)$	2	3	4	5	5	7	8	9

The 4th and 5th row of Table 1 contain  $t_g$  and  $w_g$ . The 6th row contains the needed bits for the interval  $(-w_g, +w_g)$ .

It should be remarked, that the estimation (27) and (28) are not related to optimum conditions for  $t_g$  or  $w_g$  and accordingly they give too large values.

For higher dimensions  $n$  an upper bound of  $w_{\max}$  can be used [7].

## 7. SUMMARY

Basing on the convergence proof of the error correction algorithm for ATE (5) and referring to the gap  $2\hat{\Delta}$  in  $\mathcal{X}$ -space estimations of the correction bound  $t_g$  (15) and of the sufficient interval  $(-w_g, +w_g)$  of weights (22) for digital implementation were derived. The estimation of  $t_g$  yields an upper bound for  $2\hat{\Delta}$  in dependence on the actual number of correction steps (16). In case of bounded interval of weights the estimation of  $t_g$  is preserved if the interval ensuring separation is doubled. The interval of weights ensuring separation is also given for adapting a zone  $2\hat{\Delta} < 2\bar{\Delta}$  (24). As an example the well known BOOLEAN threshold functions are used for a concrete determination of values for  $t_g$  and  $w_g$ .

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